



# Mitigating Exploits, Rootkits and Advanced Persistent Threats

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Intel Corporation

**Hot Chips Tutorial**

# Agenda

**Problem**  
**Better Protection**  
**Solid Foundations**  
**Usages**  
**Minimizing TCB**  
**Summary**

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# Increasingly Sophisticated Attacks

**Operation Aurora:** "Google announced an attack targeting it and what is believed to be more than 30 other companies ...."



*CNET January 12, 2010*

**Stuxnet:** "...a novel way to use computers to sabotage an enemy's lifeline infrastructure suggests a powerful new kind of weapon is moving within reach of weak states, militant groups and criminals..."



*Reuters Nov 30, 2010*

**The Heartbleed Vulnerability: What It Is and How It Affects You:** "...Heartbleed is not a virus, but rather a mistake written into OpenSSL" April 2014



**Banking Malware (SpyEye) Monitors Victims by Hijacking Webcams and Microphones**  
May 2012, **PCWorld**



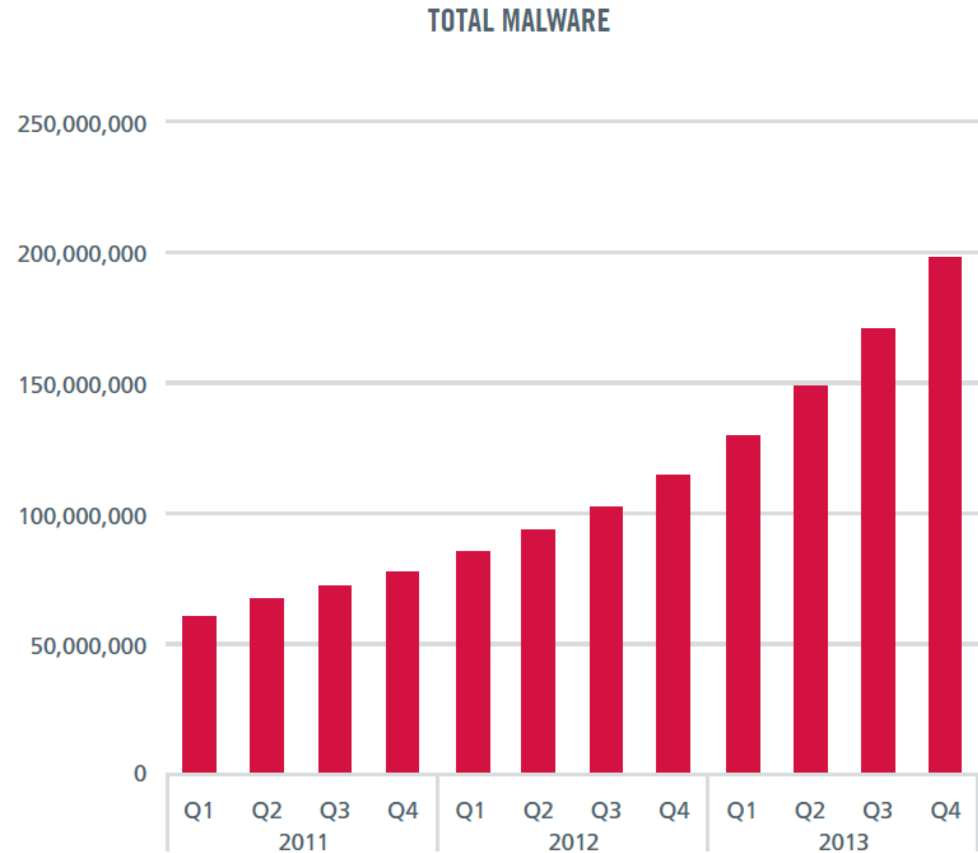
**Meet 'Flame,' The Massive Spy Malware**

**"...designed primarily to spy on the users of infected computers and steal data from them, including documents, recorded conversations and keystrokes."** May 2012, **Wired**



# Malware Signatures More & More...

- Malware samples continue sharp rise
- Polymorphic viruses
- Methods of packing, redistributing existing malware
- Looking for known malware misses 0-days and targeted attacks



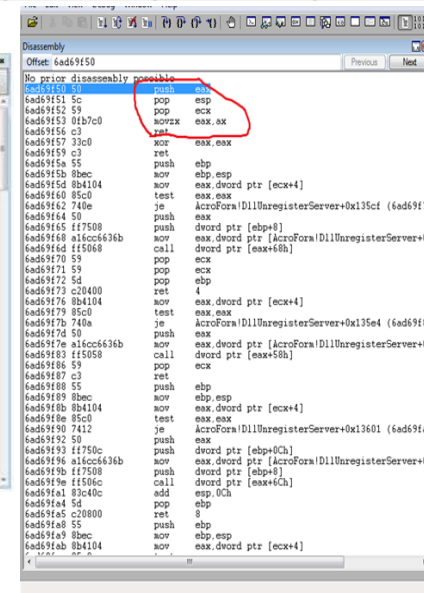
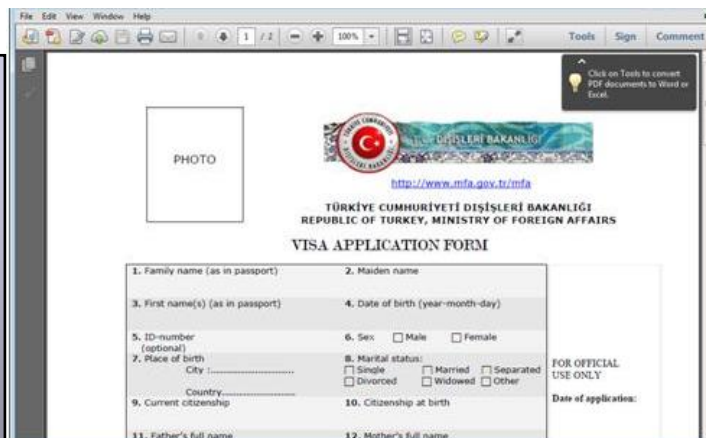
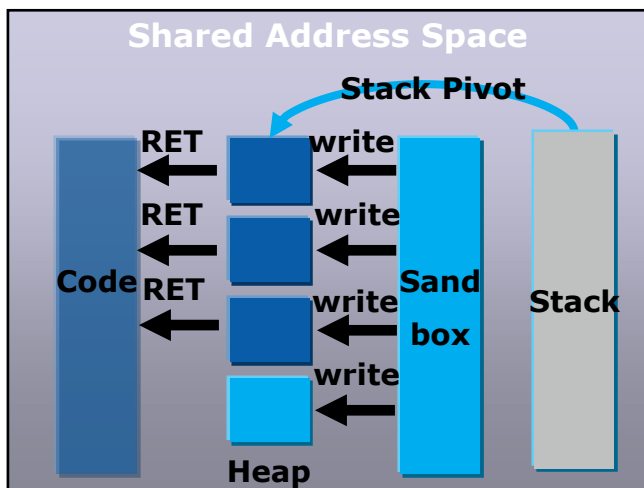
Source: McAfee Labs, 2014.

# 0-Day: Vulnerability, Armed, Exploited

## Sample: APSA13-02 Exploit Analysis

Malicious File											
Acro32		Trigger 1 <sup>st</sup> Vul	StackPivoting	D.T	2 Threads	IPC	StackPivoting	L2P.T	L2P.T	L2P.T	Acro32
Sandboxed Reader	Open PDF file	Heap Overflow	Run ROP Stackpivoting	Create and Load D.T Library	1: Show Error 2: Create L2P.T	IPC to trigger 2nd vul, then quit					New Process shows Visaform Turkey.pdf
Normal Reader						Heap Overflow	Run ROP via Stackpivoting	Load L2P.T	Create Langbar.dll And load it	Create Visaform Turkey.pdf	Create new sandboxed process to open new pdf

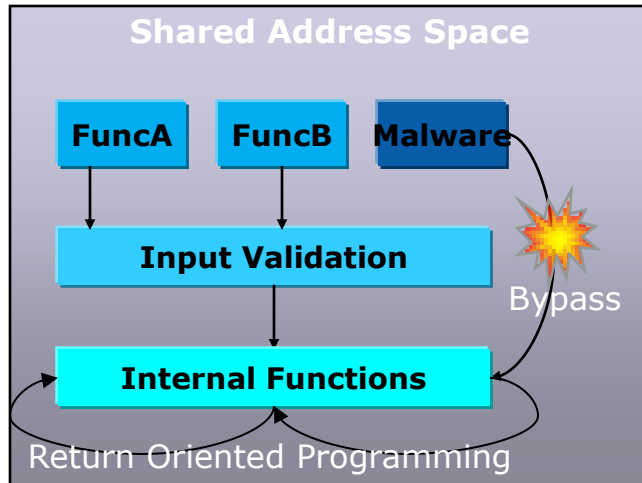
## Stack Pivot



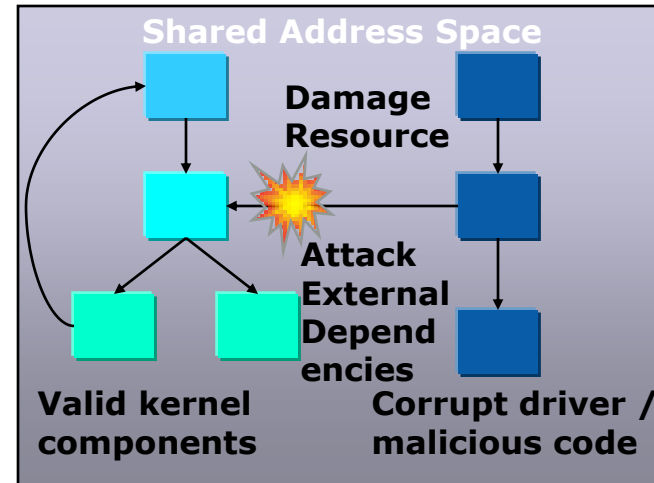
## Emerging 'Stack Pivoting' Exploits Bypass Common Security- APSA13-02 exploit:

# Generalized Attack Vectors

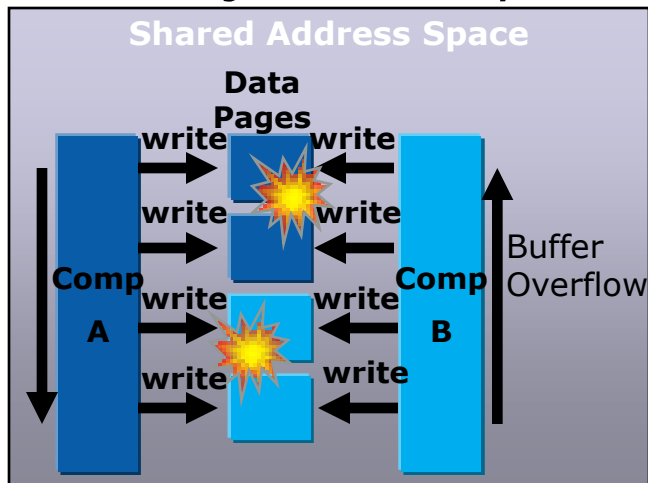
## Circumvent



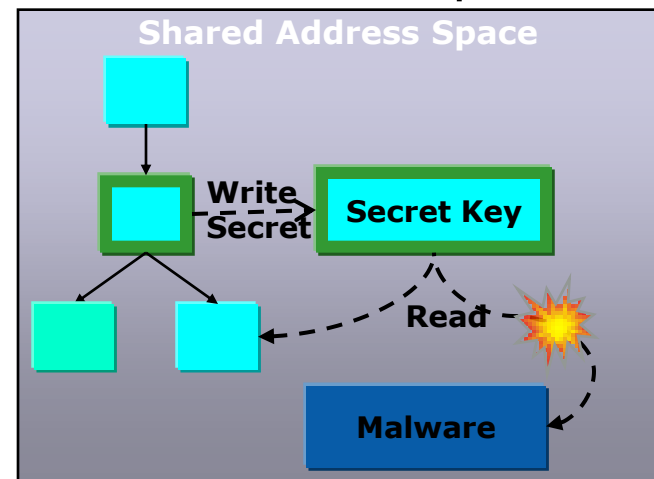
## Disable



## Inject/Modify



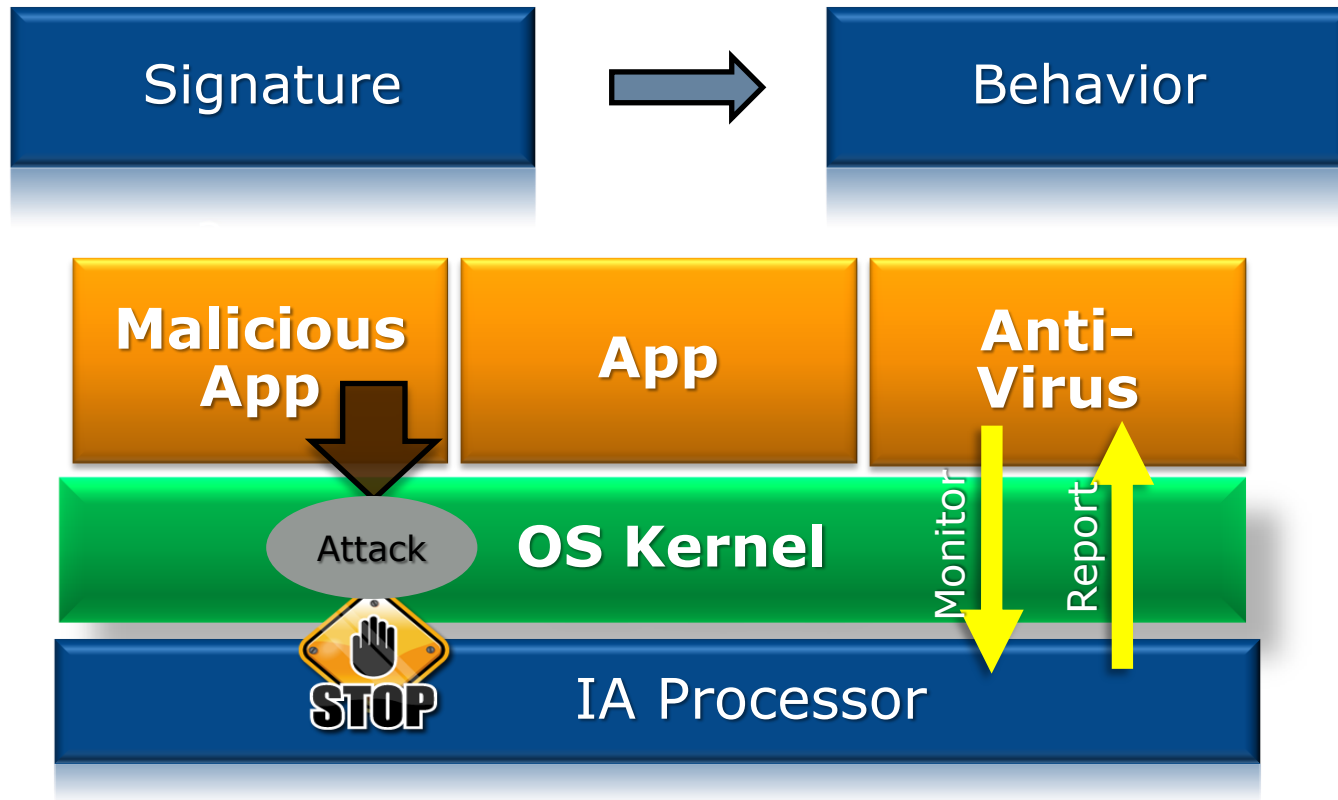
## Eavesdrop



**Problem**  
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# Paradigm Shift



Monitor behavior

Detect onset of the attack

Prevent suspicious access

*“Behavior-based” detection to stop zero-day attacks*

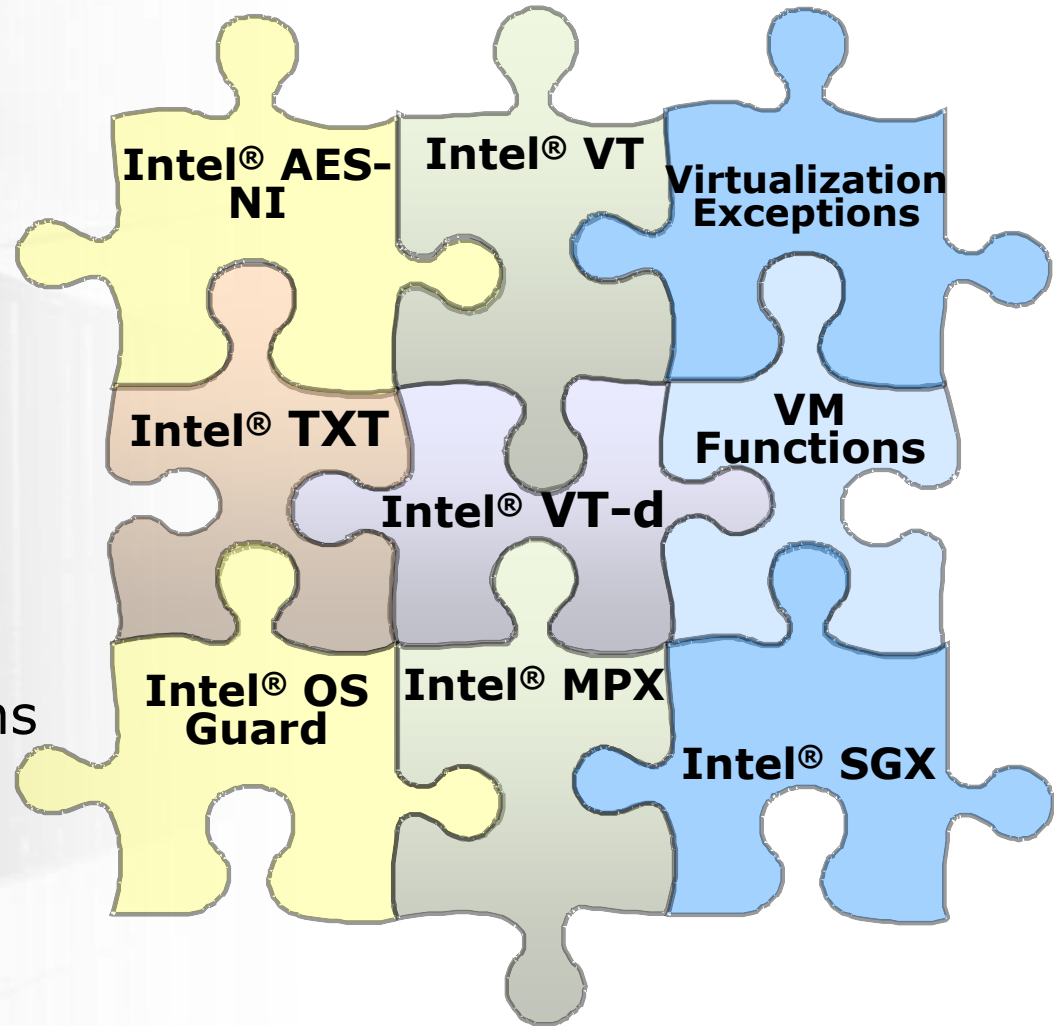
# Putting the Defenses Together

- Trusted Launch
- Measurement
- Software Isolation
- Device Isolation
- Crypto Acceleration
- Introspection Acceleration
- Supervisory Execution Prevention

Forward Looking...

- Buffer Overflow Protections
- Minimal TCB

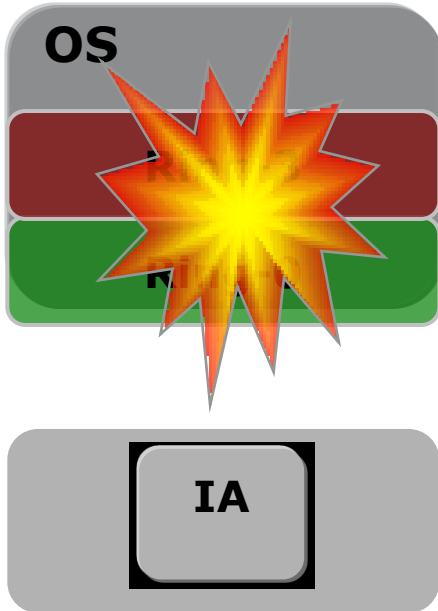
...



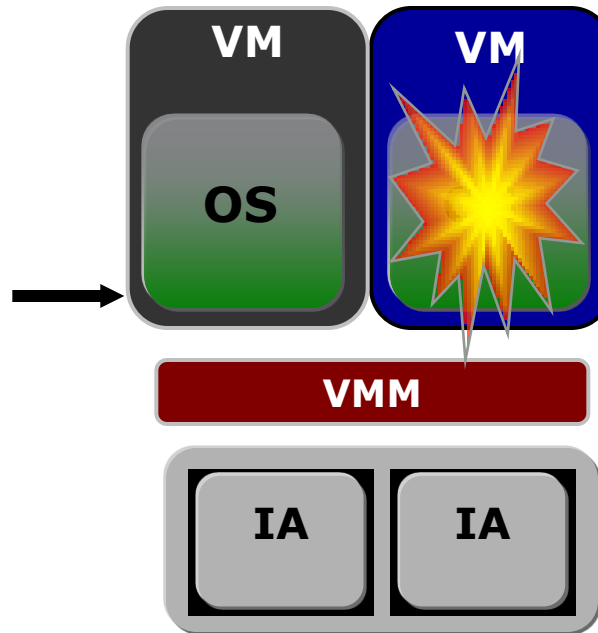
Intel® Trusted Execution Technology (Intel® TXT)  
Intel® Virtualization Technology (Intel® VT) for Directed I/O (Intel® VT-d)  
Intel® Virtualization Technology (Intel® VT)  
Intel® Memory Protection Extensions (Intel® MPX)  
Intel® Software Guard Extensions (Intel® SGX)

# Evolution of Memory Protections

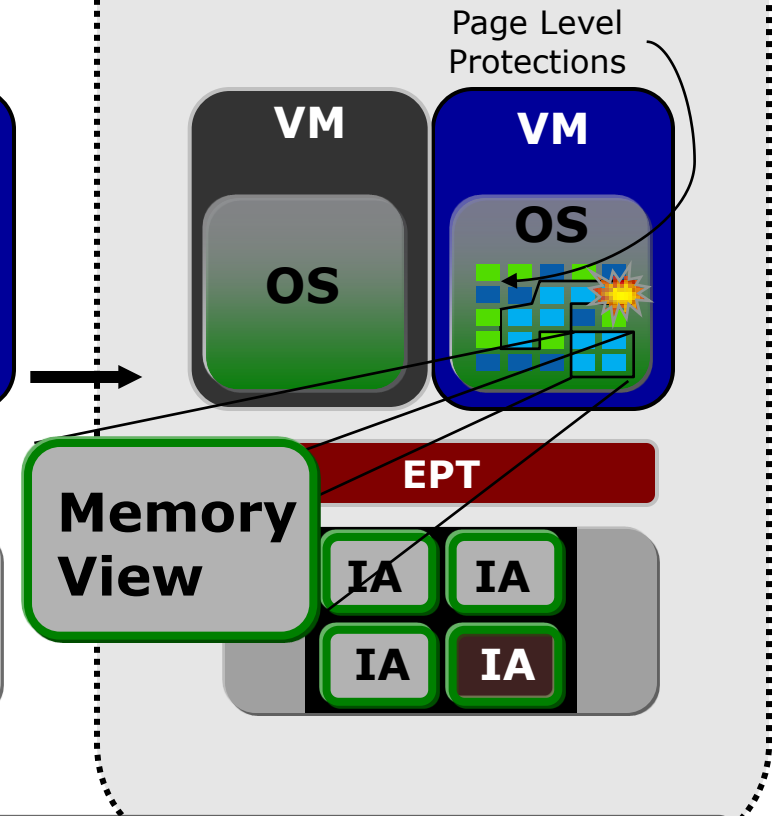
**Privilege levels based on protection rings**



**Multiple virtual machine isolation**

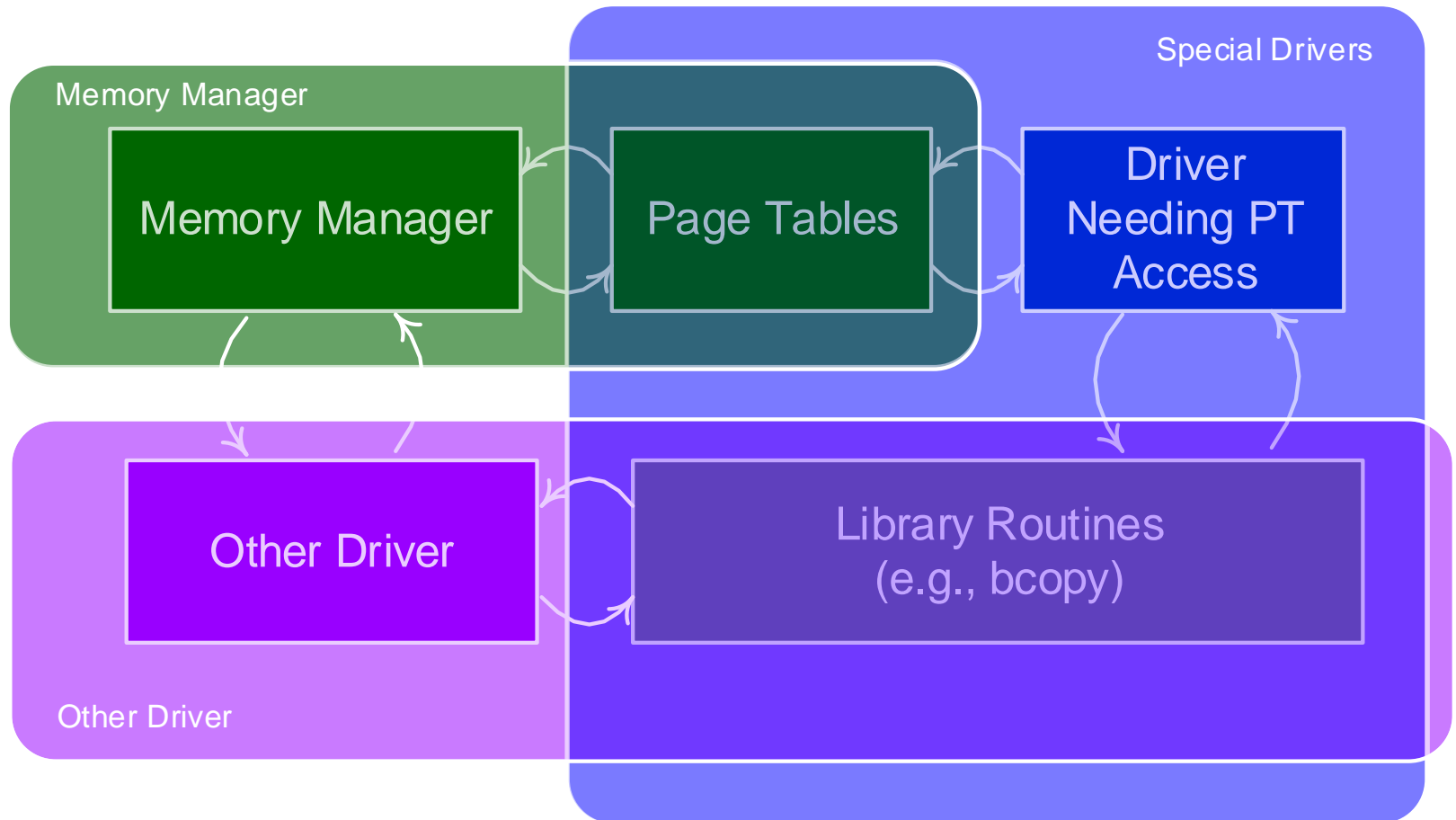


**Granular isolation within an address space**



***Reduce the attack surface while minimizing overhead***

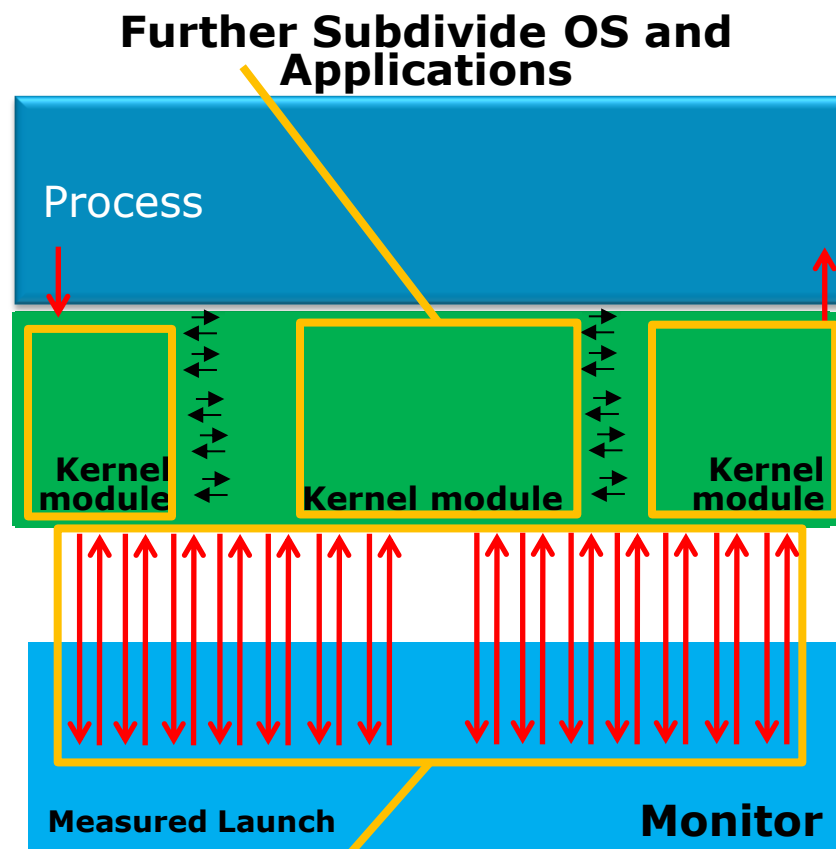
# Overlaying Granular Protections



***Accommodates existing OS methodology and legacy code***

# Accelerating Granular Isolation

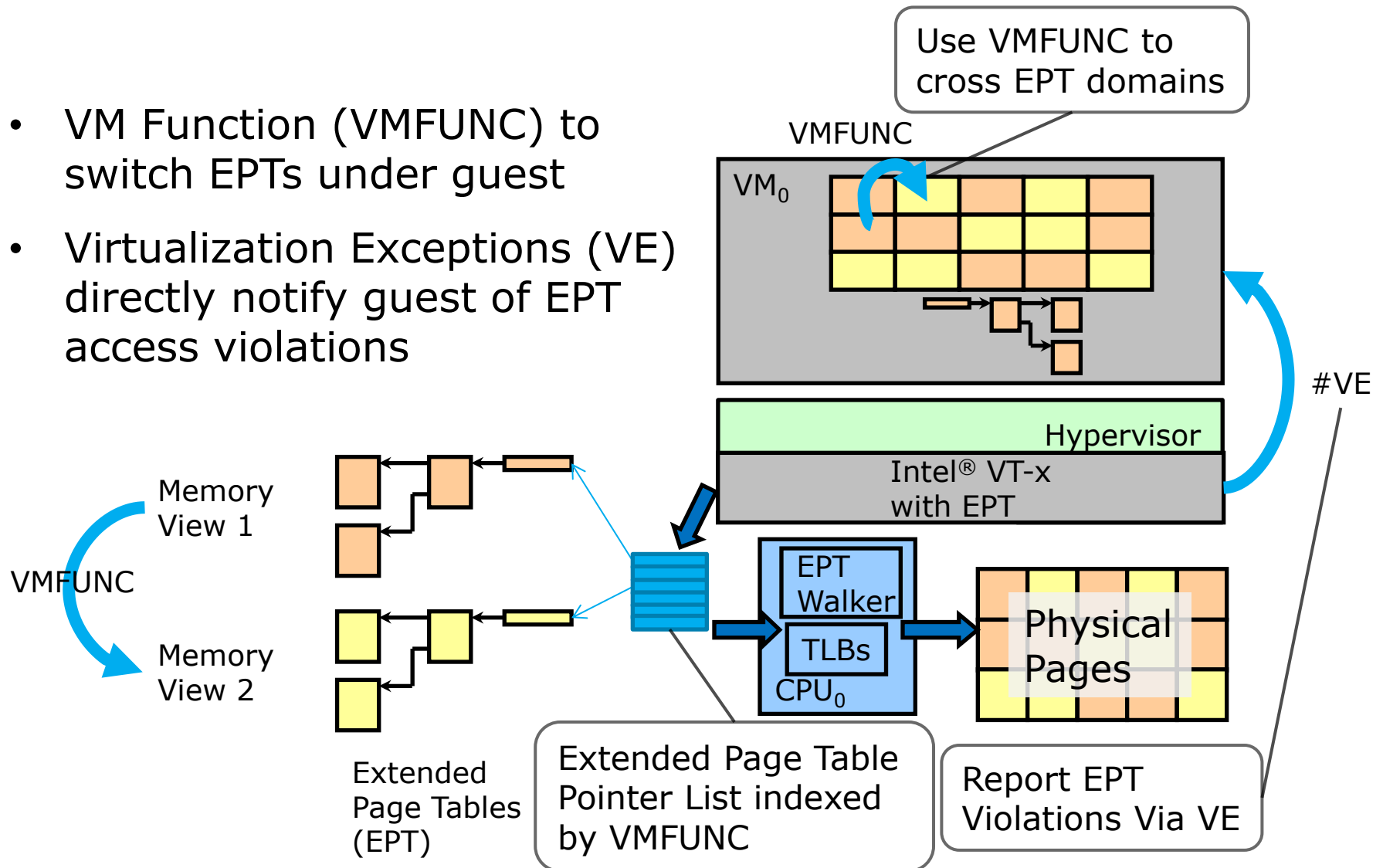
- Intel® Virtualization Technology (Intel® VT) enables protections beyond the OS
- Overlays additional protections and monitoring policies by enabling memory views
- Provides continuous detection of illicit behaviors
- Accelerated using VM Functions and Virtualization Exceptions...



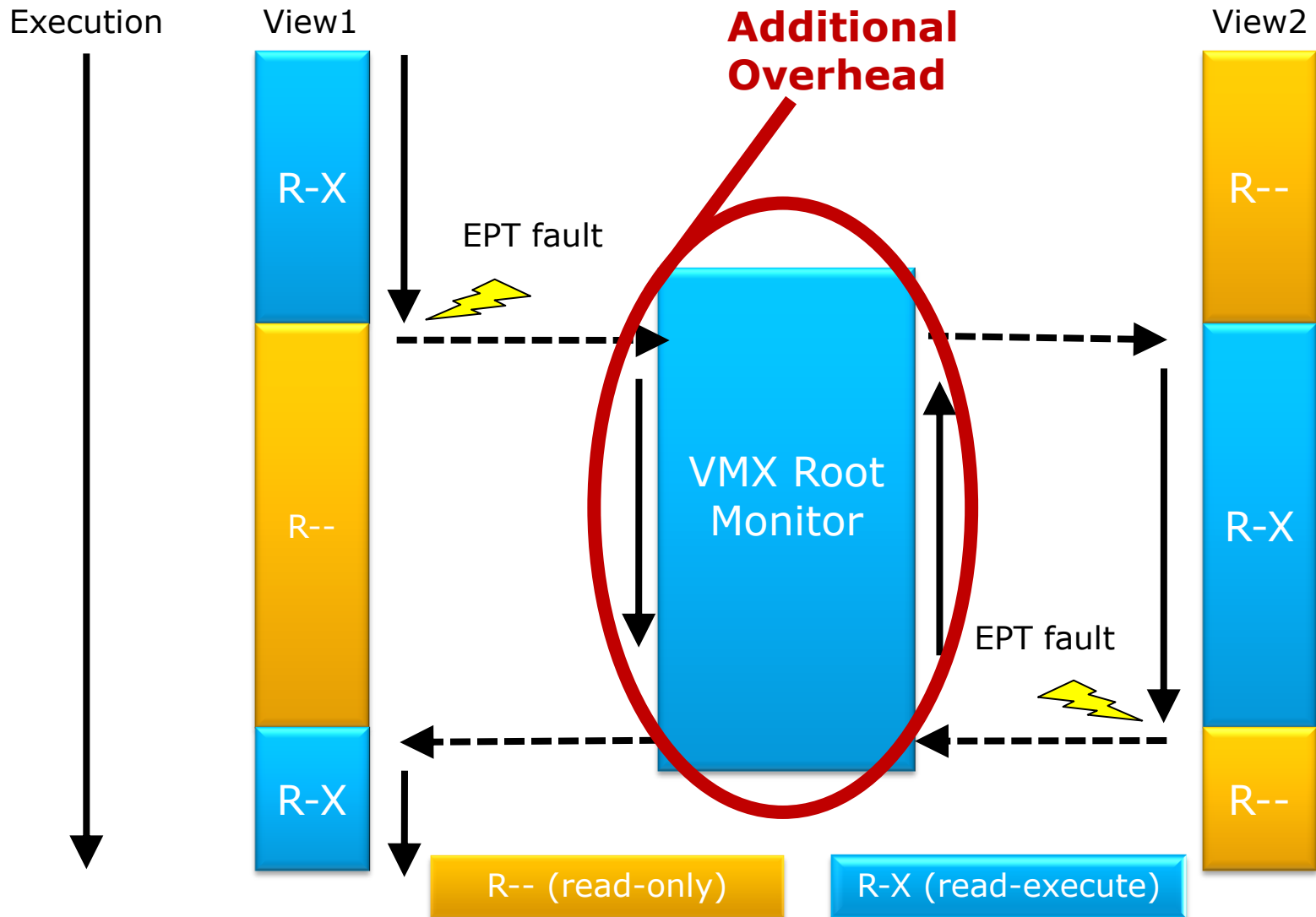
**Additional Monitoring:** Privileged software monitors OS activity

# Extended Page Tables for Isolation within VM

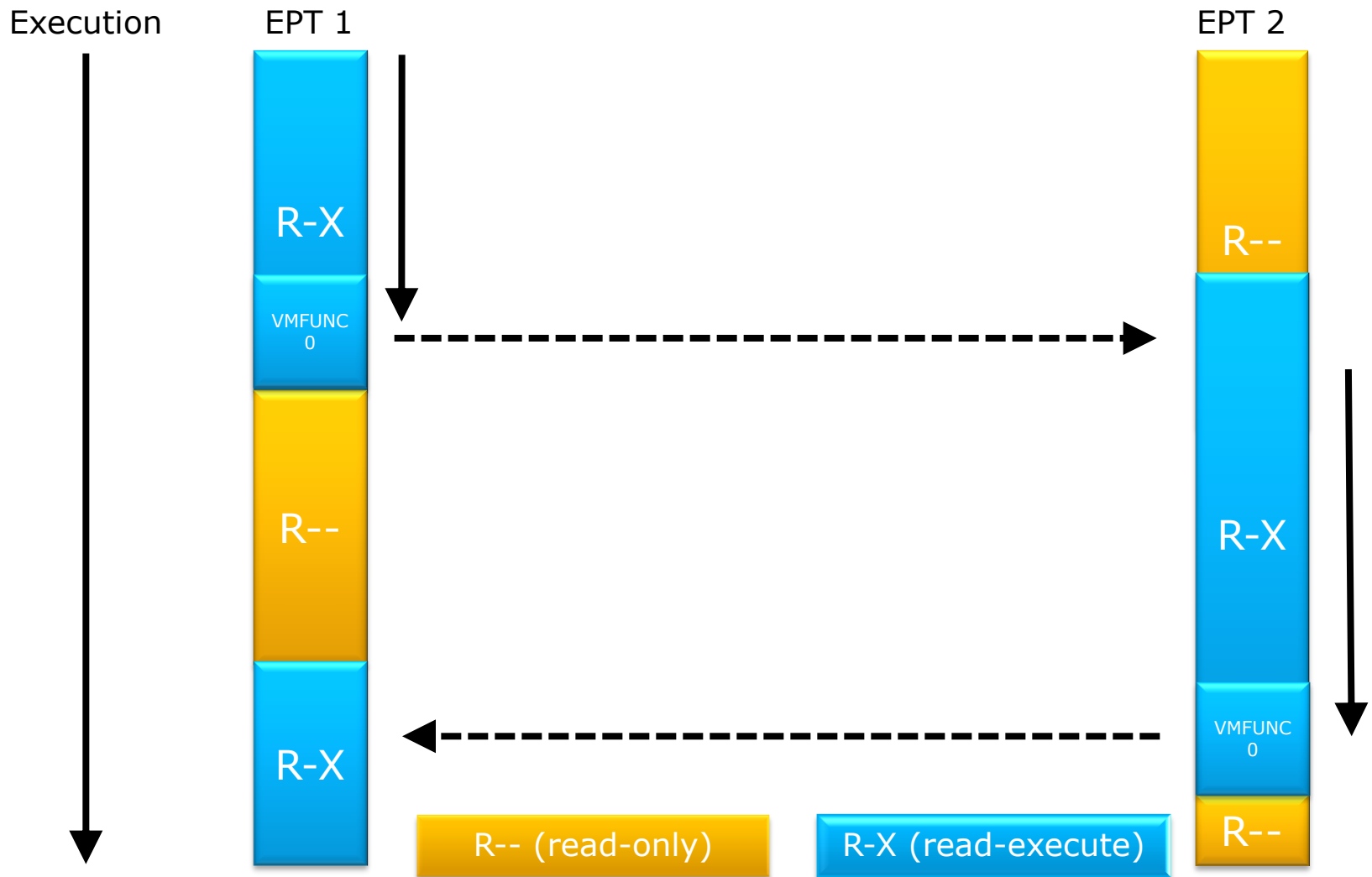
- VM Function (VMFUNC) to switch EPTs under guest
- Virtualization Exceptions (VE) directly notify guest of EPT access violations



# Granular Isolation (Before)



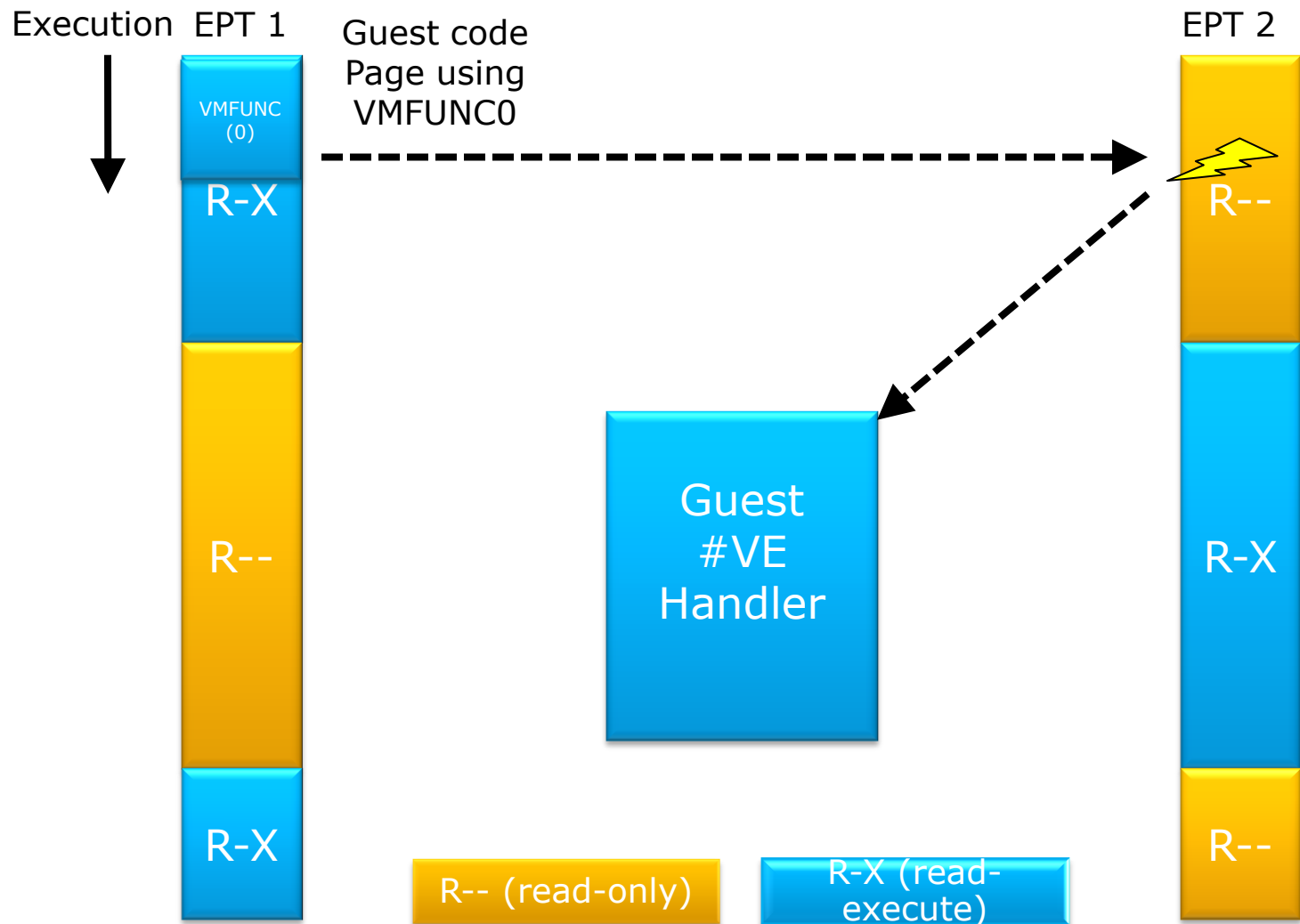
# Granular Isolation with VM Functions



***Enforce Control Flow Integrity with Intel® VT***



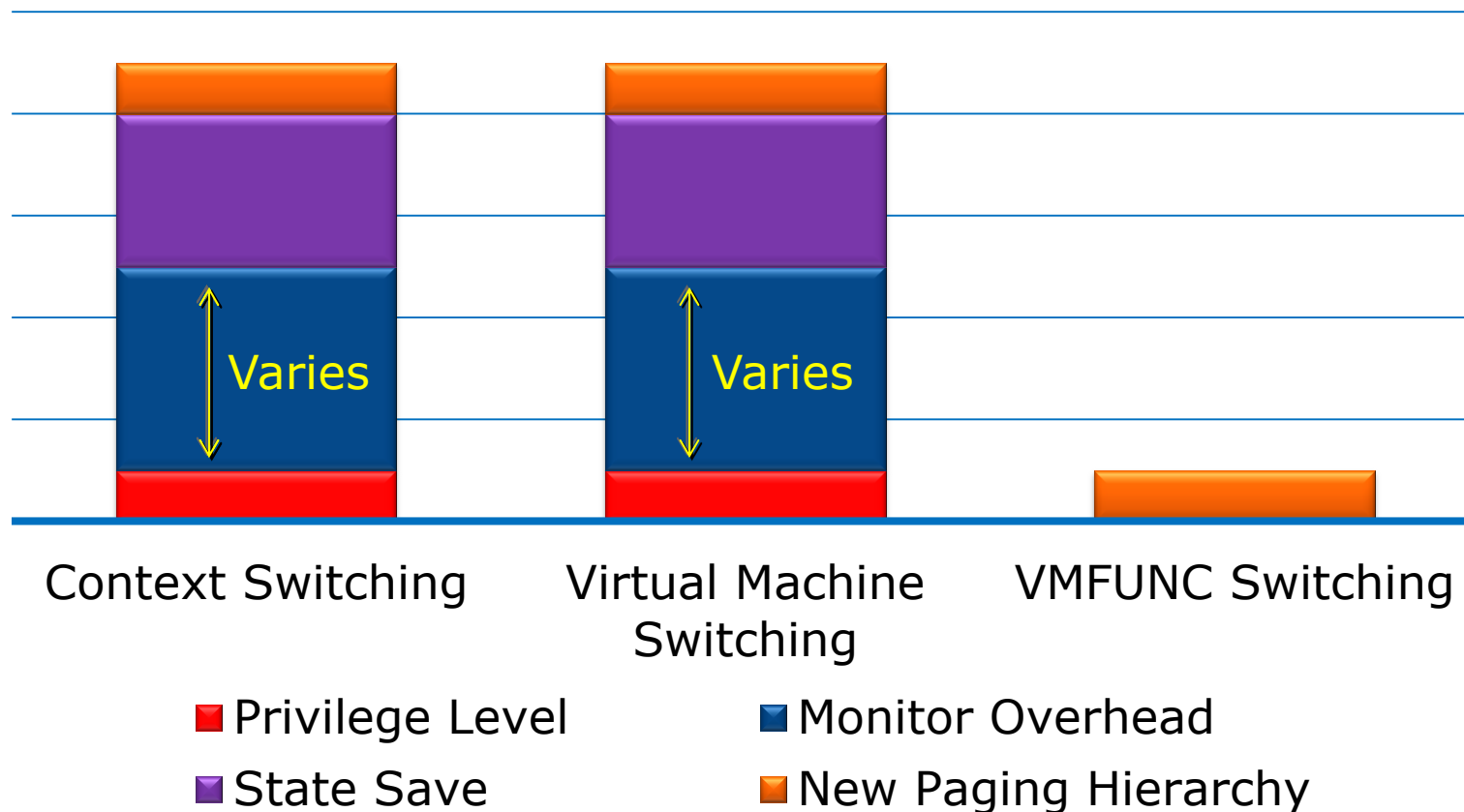
# Monitoring with Virtualization Exceptions



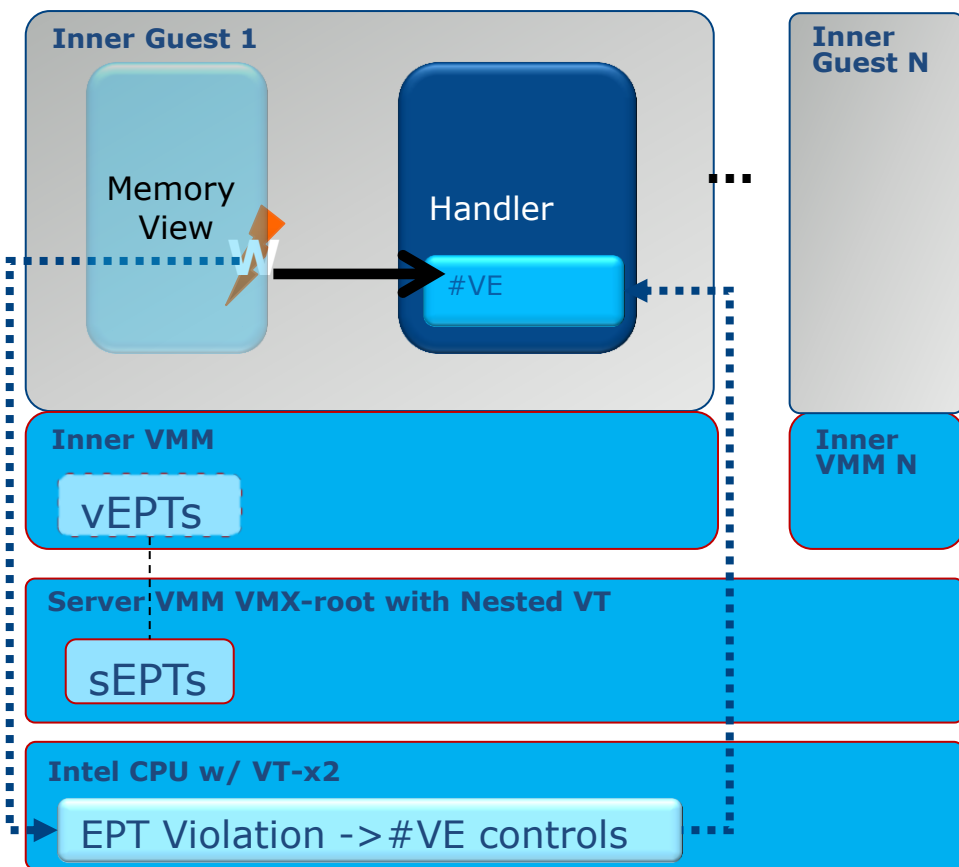
***EPT Exceptions directed to the guest***

# VM Function for Switching

## Relative Performance Comparison



# Layering Virtualization and Introspection

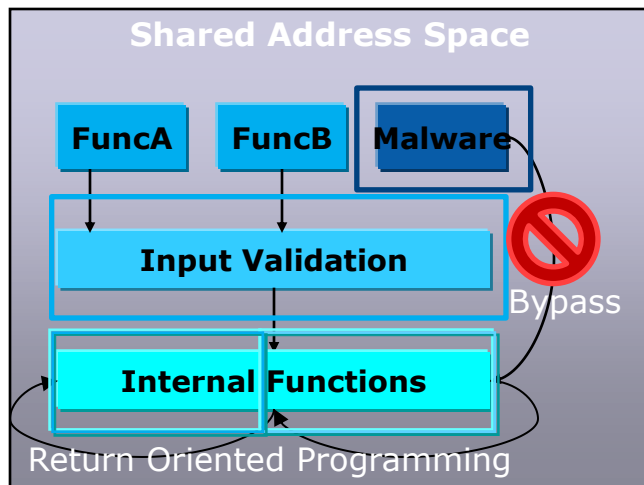


- Root VMM supports VT nesting (with EPT shadowing)
- Root VMM and guest opt-in to enable features
- VMFUNC switches authorized EPTs without engaging VMM(s)
- EPT violations reported via #VE to guest directly
  - No VM Exits for guest policies
  - No additional overhead for VMMs
- Root VMM decides which pages #VE and which will VM Exit
  - Disambiguates copy-on-write and other VMM notification needs

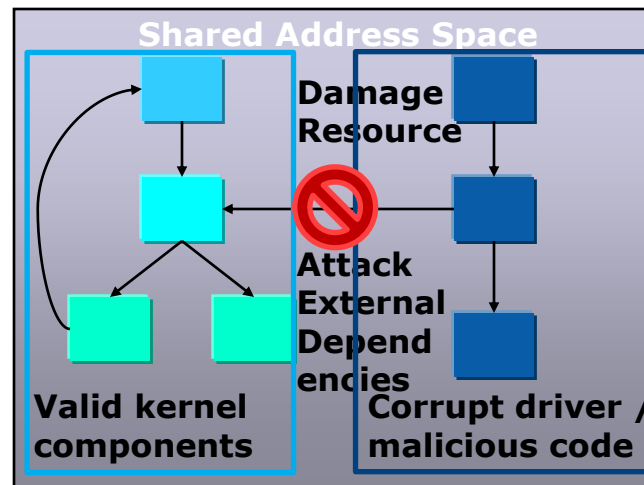
***Enables efficient introspection across multiple VMs***

# Revisiting Attack Vectors

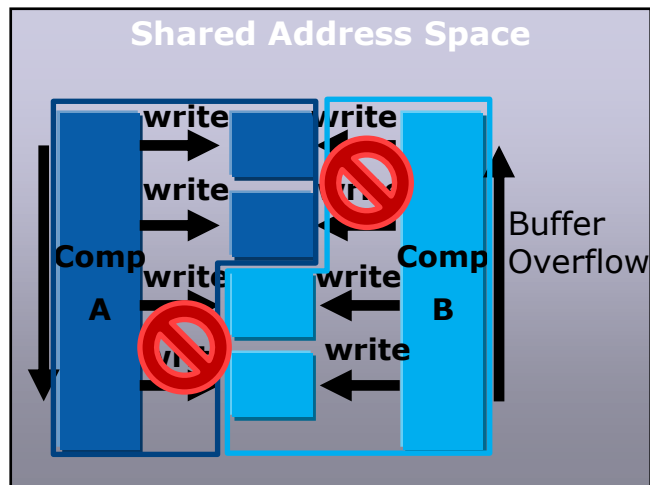
## Circumvent



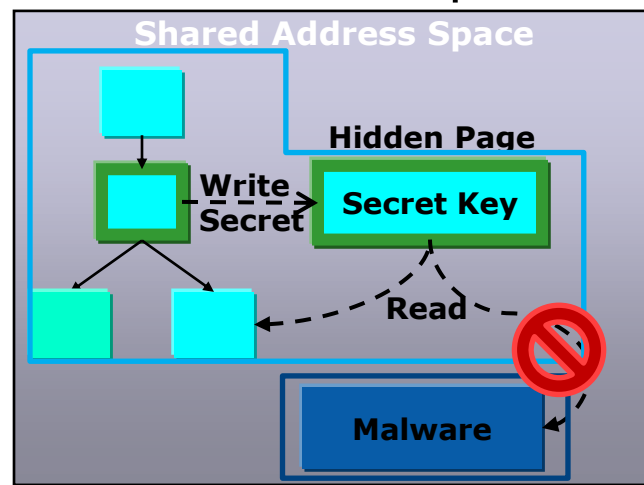
## Disable



## Inject/Modify



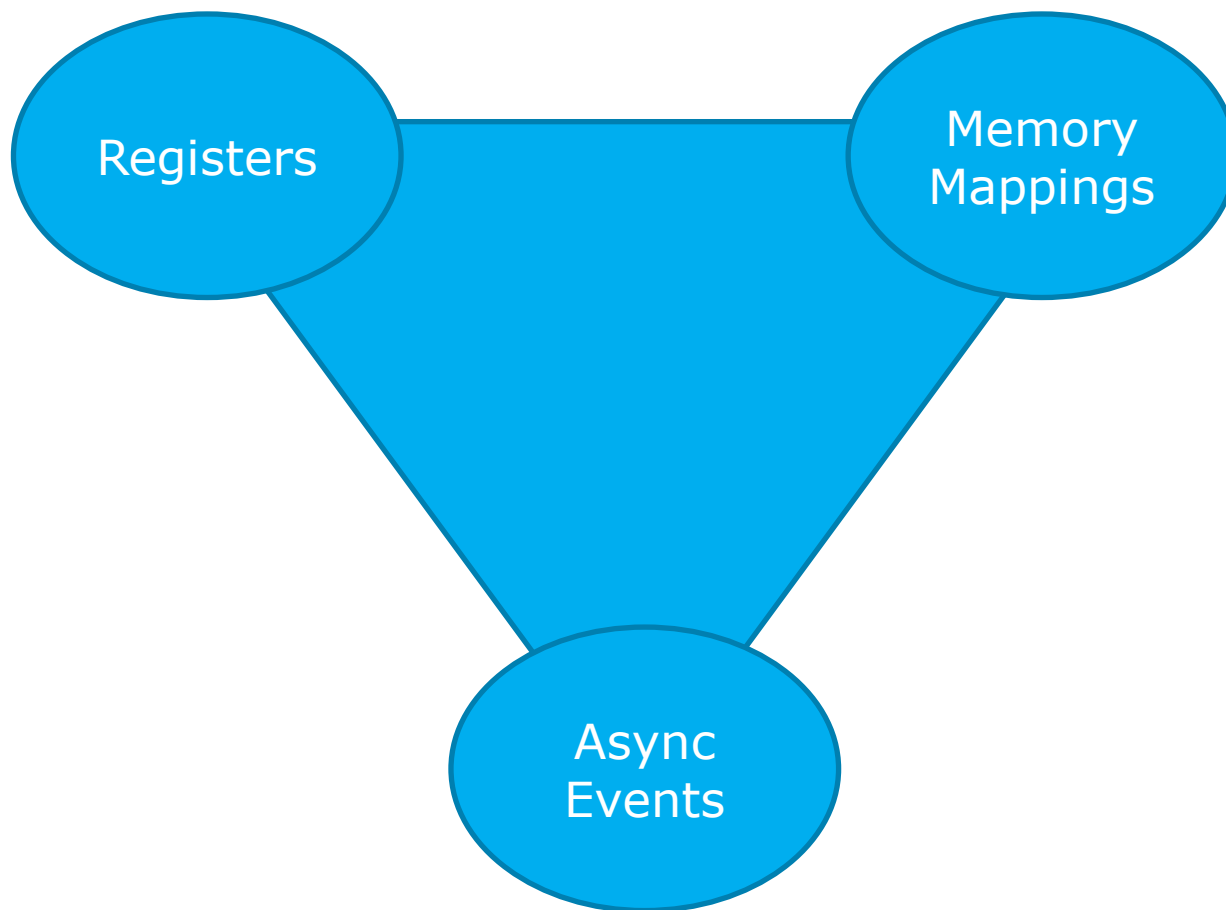
## Eavesdrop



***Overlay memory views to monitor software behavior***

**Problem**  
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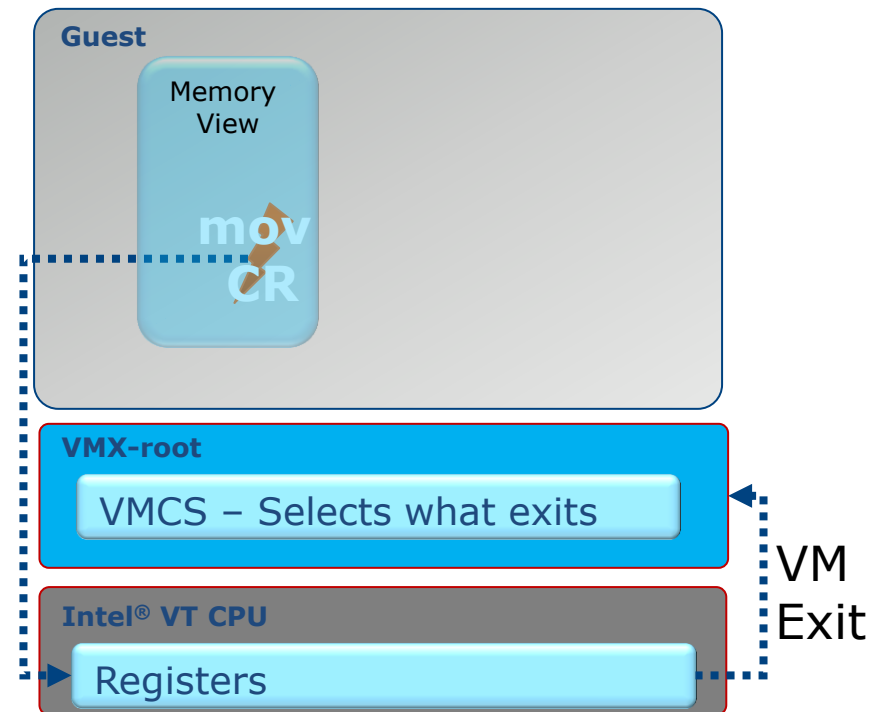
# Overlaying Protections: Things to Consider...



*May also address external devices depending on usage*

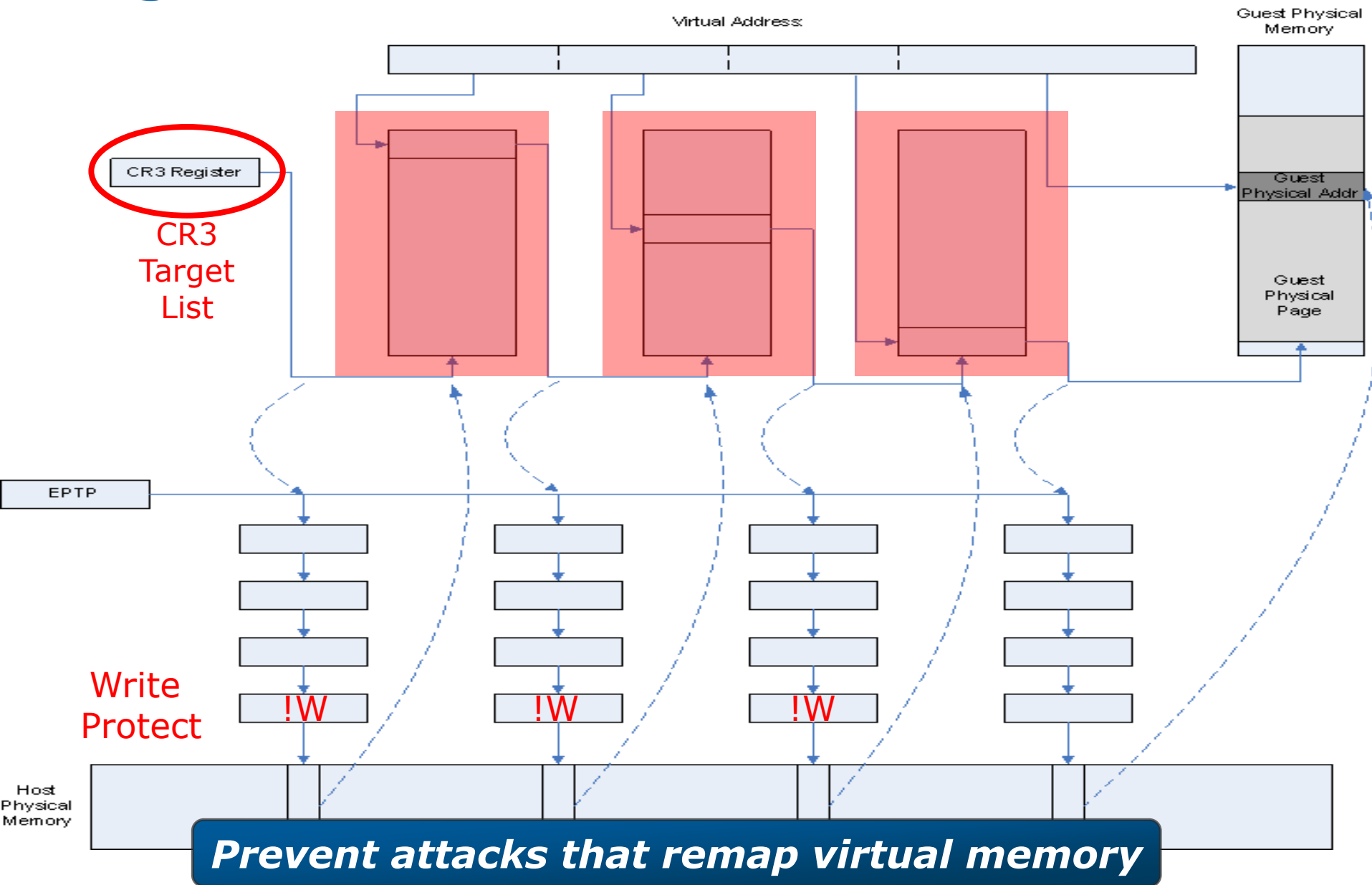
# Monitoring Processor Registers

- VMM can be configured to intercept changes to:
  - Model Specific Registers
  - Control Registers
  - Debug Registers
  - Descriptor Tables (IDTR...)
  - Mode dependent...
- VMCS determines what registers to monitor
- GPRs including the stack pointer can be checked at boundaries and on events



***Monitor processor state to prevent attacks***

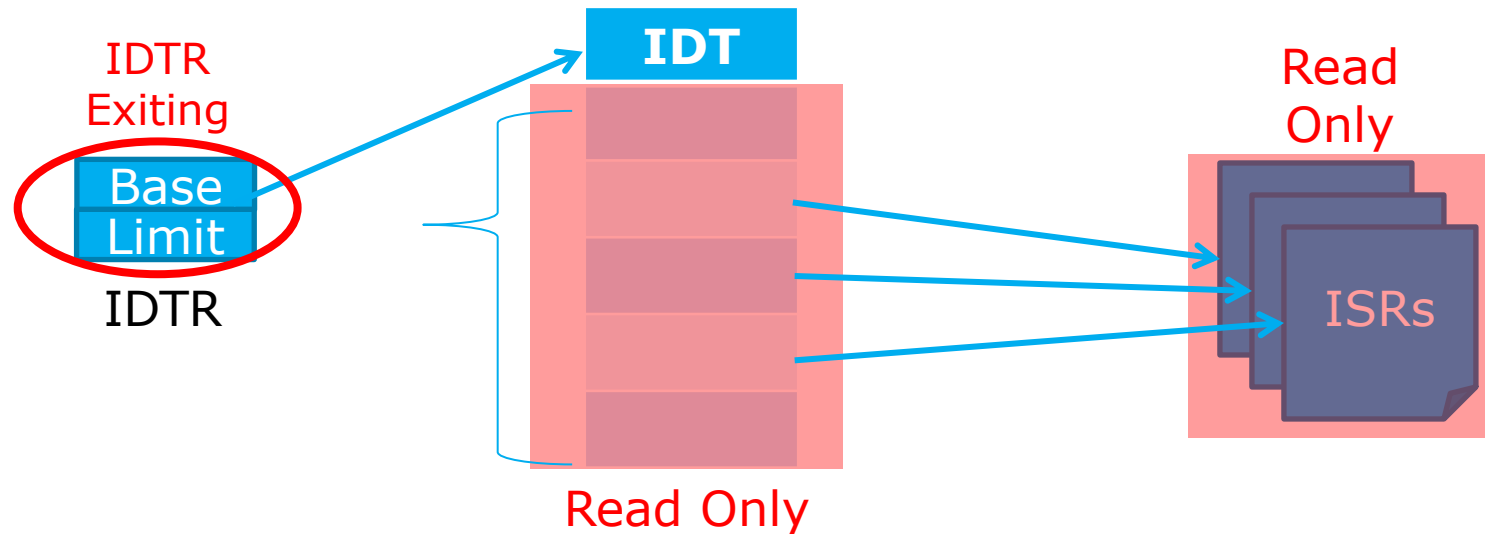
# Page Table Edit Control





# Interrupts & Asynchronous Events

- Protect Interrupt Descriptor Table & Register
- Trust Interrupt Service Routines or own ISR stub
- Stub code protects state
  - Stack
  - General Purpose Registers

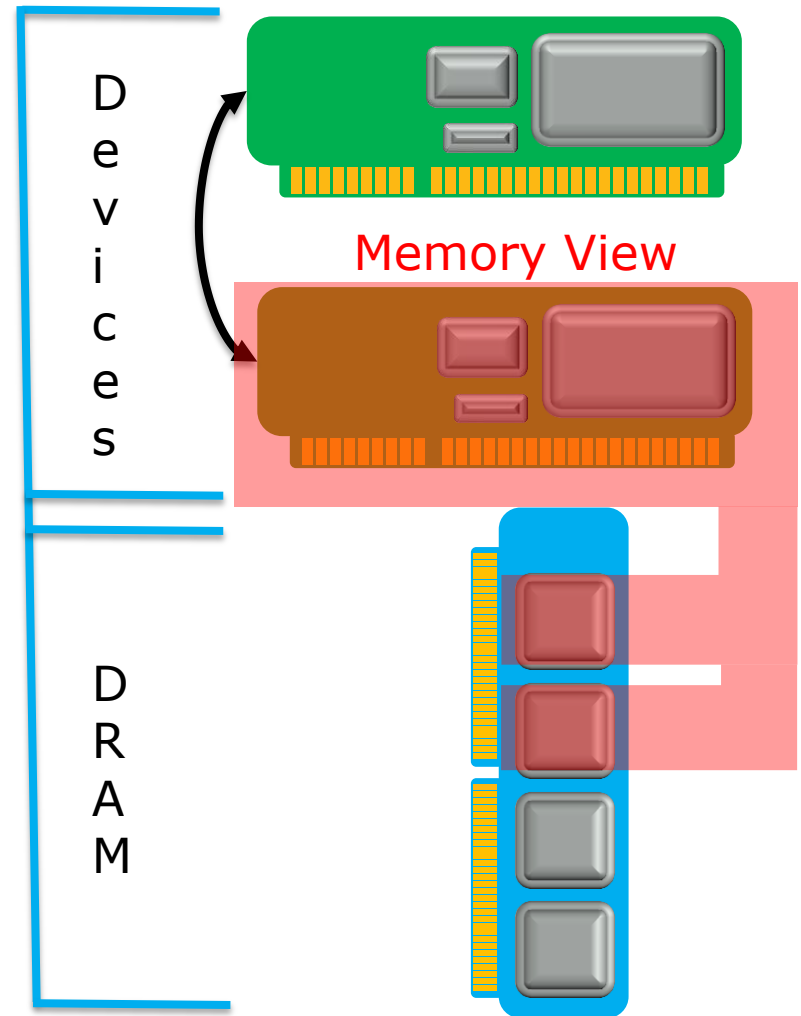


***Protect or intercept asynchronous paths***

# Devices

- Device space configuration
  - Programmed IO in/out
  - Relocation of device registers in memory/BAR change
  - Trigger VM Exit
- Memory Mapped IO
  - Device registers
  - Covered by EPT policy
- DMA
  - Buffers in memory
  - Covered by EPT policy

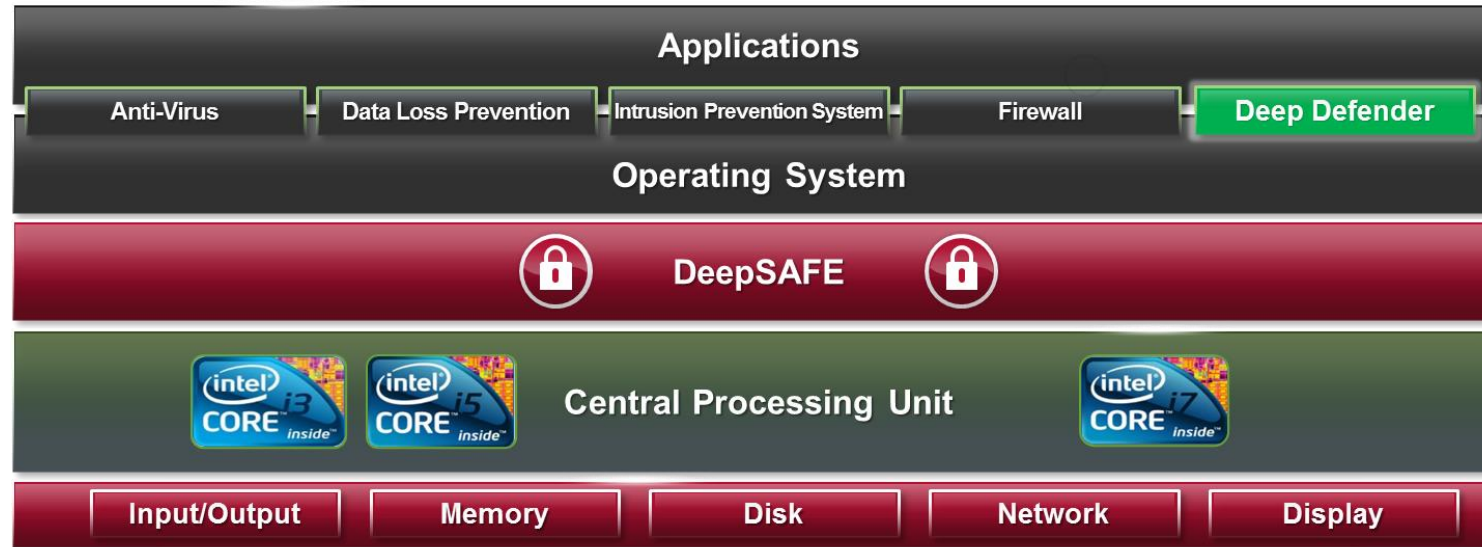
Address Space



***Intel® VT-d protects against compromised devices***

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# McAfee® Deep Defender Overview



- McAfee® DeepSAFE™ technology in the McAfee® Deep Defender product can safely monitor writes to critical memory assets
- The Deep Defender component within the operating system understands the O/S layout and rootkit techniques
- The DeepSAFE component uses CPU primitives to monitor CPU and memory so that pages containing sensitive code and data are access-controlled

# Providing Better Protections



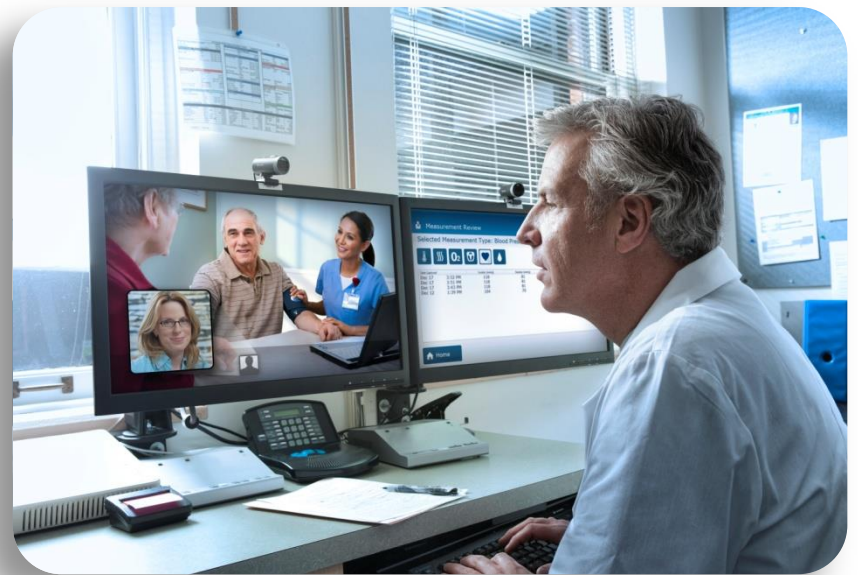
Input



Audio



Storage

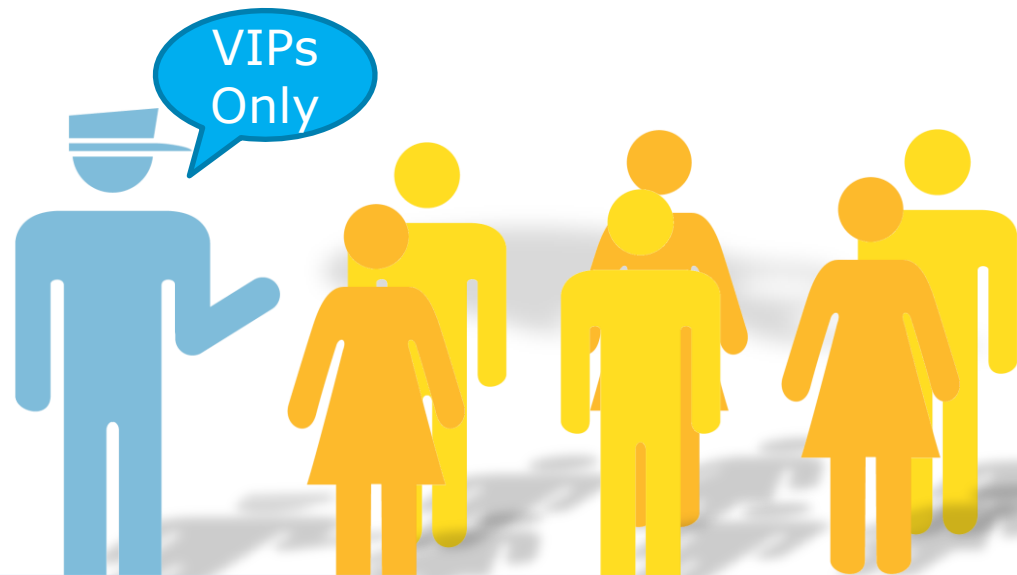


Video  
Hot Chips 2014 | Tutorial

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# Democratizing Security

- Software relies on a Trusted Execution Environment (TEE) in case other defenses fail
- Isolated hardware and restricted modes limit use
- A future where there are enough TEEs for all?

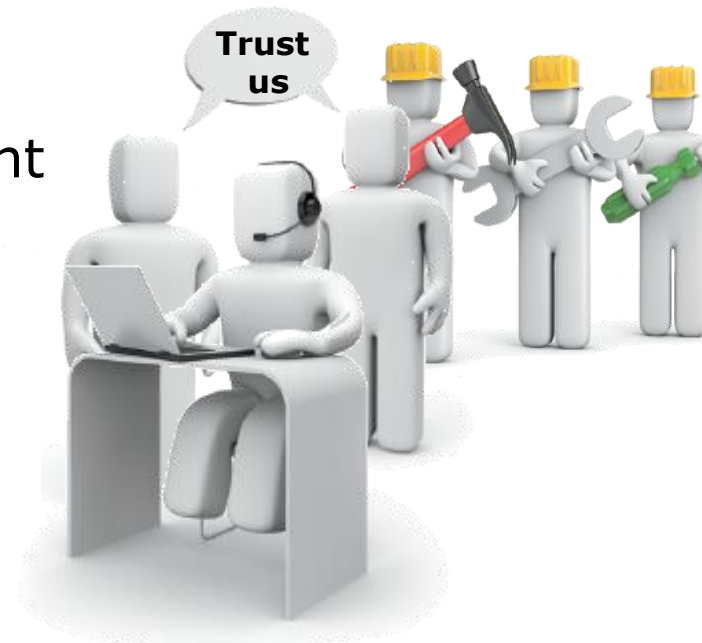


*Scaling Trusted Execution Environments for the many*



# Trust How Many?

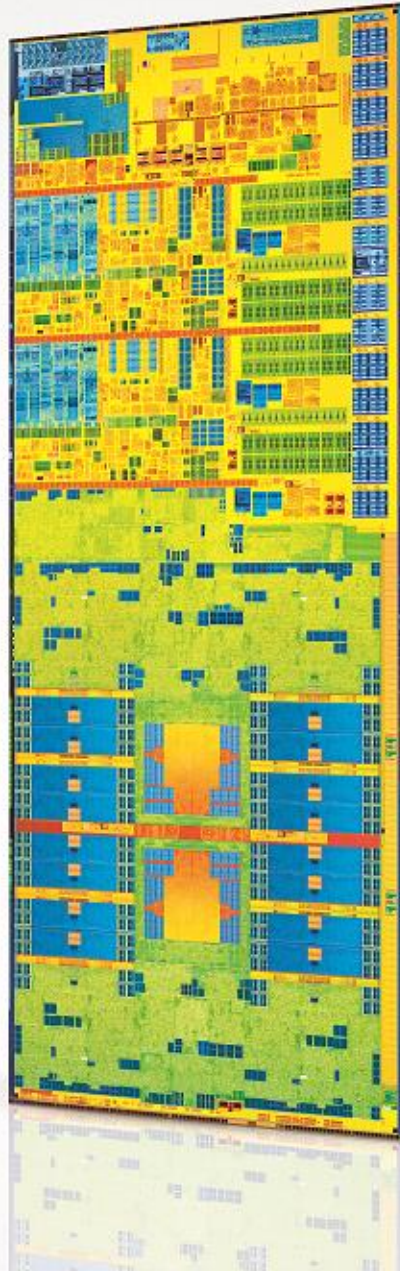
- Trusted Computing Base (TCB):
  - Set of all hardware, firmware & software part of a trusted environment
- Minimizing the TCB:
  - Remove Software Stacks
  - Remove Drivers
  - Remove Devices
  - Remove Firmware...



# TCB



# Trust The Processor

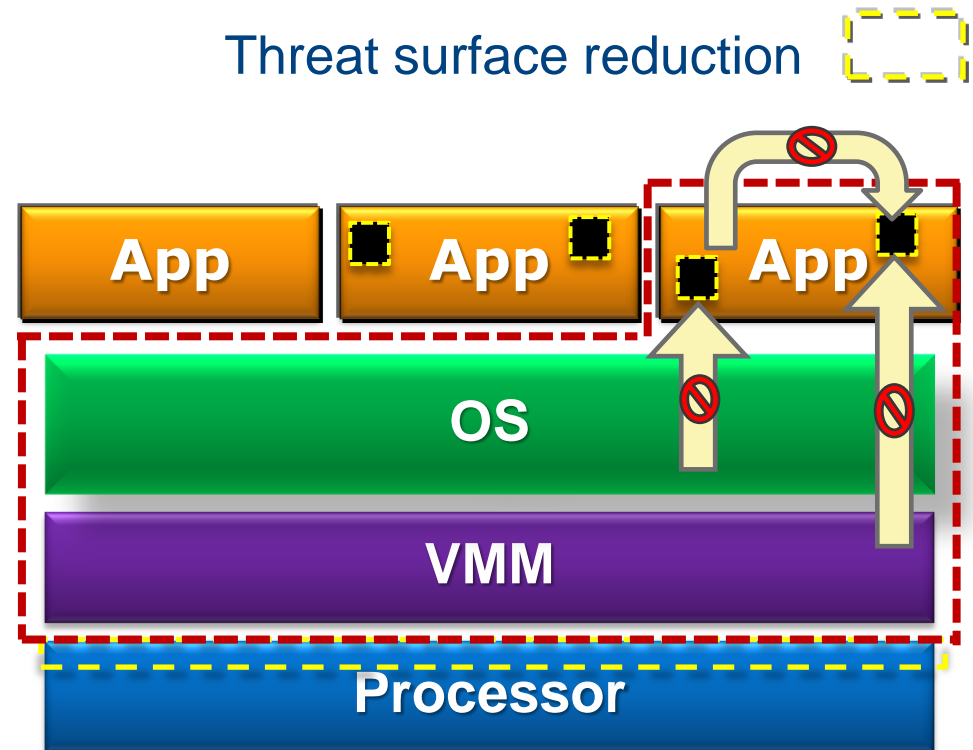


# Scaling Trust with a Minimal TCB

Define precise trust boundaries

Remove extraneous components from the trust boundary

Reduce TCB to the processor



***Scalable security within mainstream environment***

Intel® Software Guard Extensions Programming Reference:  
<https://software.intel.com/sites/default/files/329298-001.pdf>

# Utilize Existing Instruction Set Security Primitives

E.g. Intel® Advanced Encryption Standard New Instructions:

AESKEYGENASSIST

*ShiftRows()*  
*SubBytes()*  
*MixColumns()*  
*AddRoundKey()*

AESIMC

*InvShiftRows()*  
*InvSubBytes()*  
*AddRoundKey()*

AESENC

*InvShiftRows()*  
*InvSubBytes()*  
*InvMixColumns()*  
*AddRoundKey()*

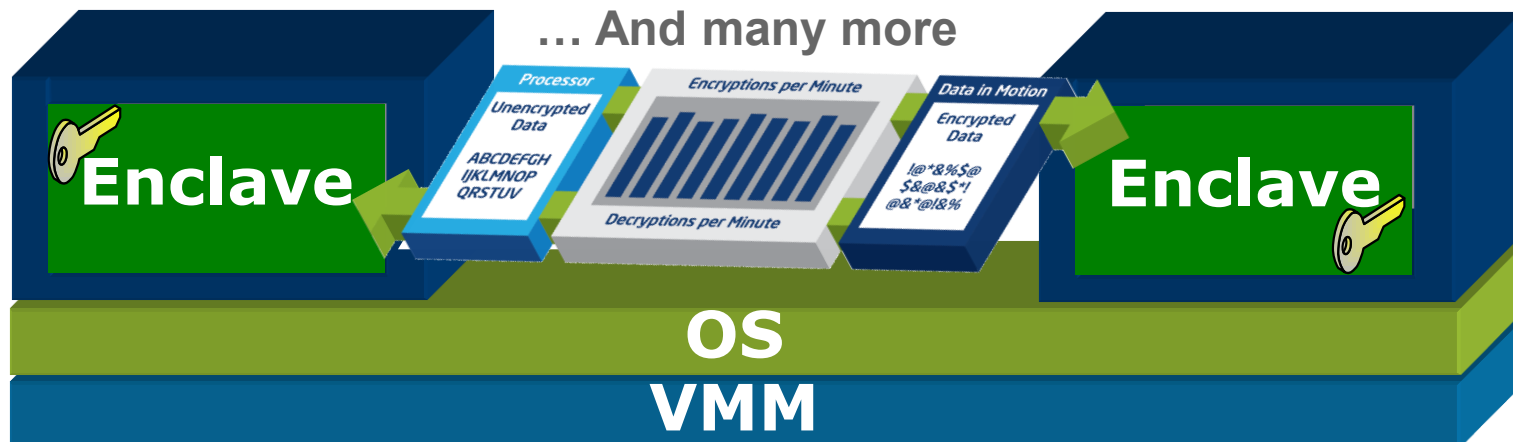
AESENCLAST

*ShiftRows()*  
*SubBytes()*  
*AddRoundKey()*

AESDEC

AESDECLAST

... And many more



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# Summary

- Increasingly sophisticated attacks require better defenses
- Moving from signatures to behavioral models
- Next generation processors deliver new capabilities for advanced software monitoring and protection
- Ability to layer protections over legacy software
- Minimizing the Trusted Computing Base is the next step...

# Biography



David Durham is a Senior Principal Engineer and Director in Intel Labs. His research team developed anti-malware and cryptographic security features currently found in hundreds of millions of Intel processors. David also developed policy-based network management technologies, created security solutions shipping in Intel® vPro™ platforms and worked with McAfee to deliver virtualization-based anti-malware products. Collaborating with industry leaders, his team developed IEEE 802.1 security protocols and advanced network access control capabilities now embedded in tens of millions of Intel platforms. He is a prolific author on computer communications, having written a book, multiple publications and several Internet protocol standards deployed in millions of connected devices. David received two Intel Achievement Awards, was granted over 100 US and international patents and earned his B.S. and M.S. degrees in Computer Engineering from Rensselaer Polytechnic Institute.

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